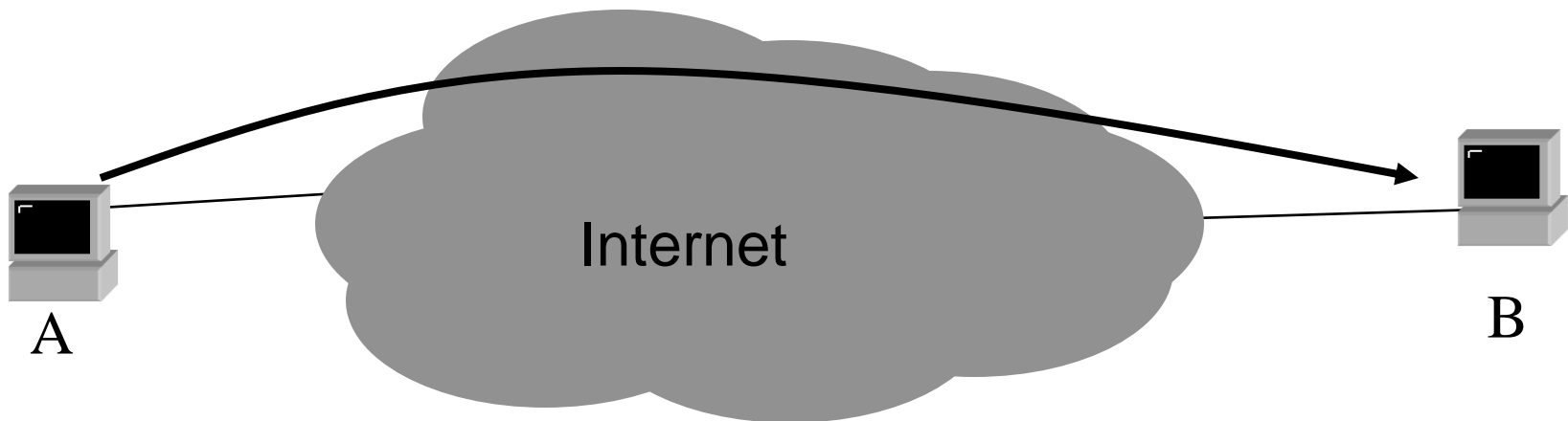


Introduction to IP Routing

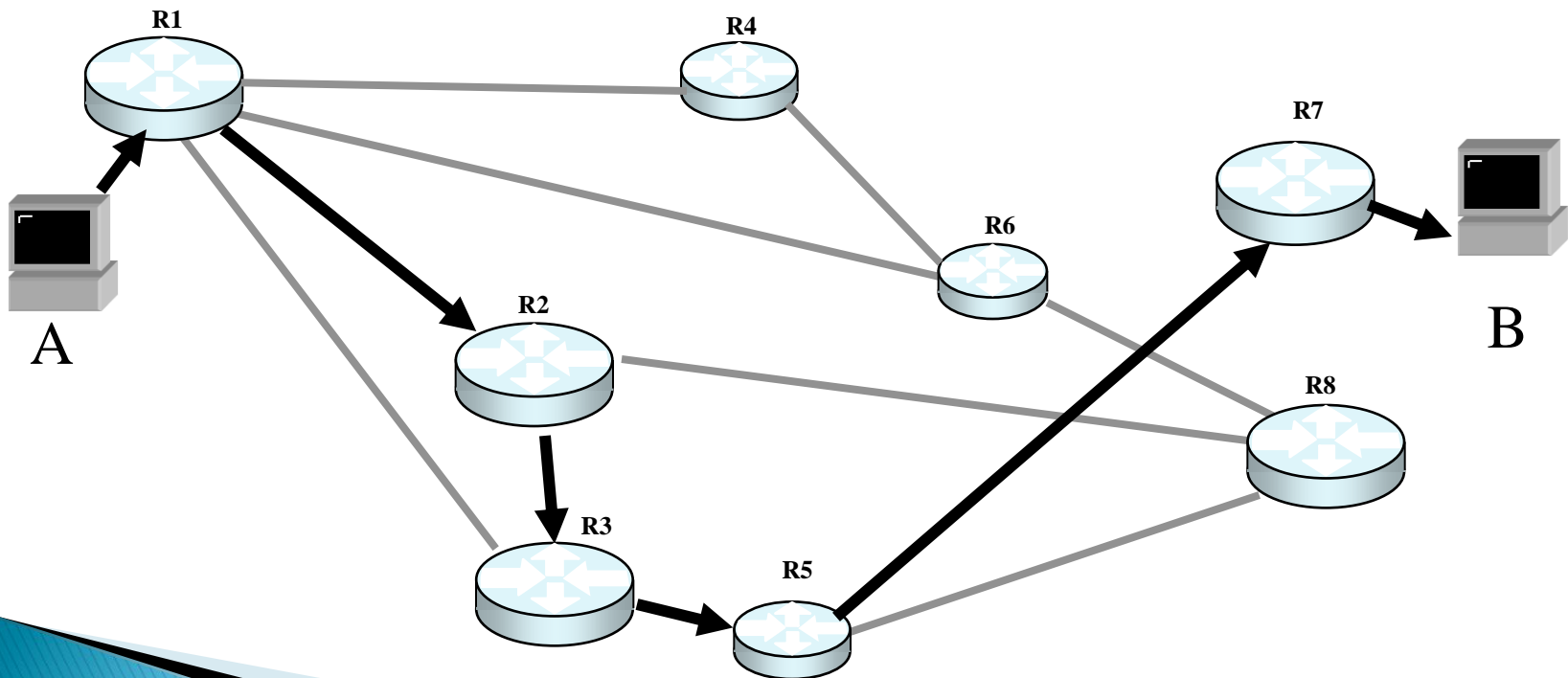
Routing

- ▶ How do packets get from A to B in the Internet?



Connectionless Forwarding

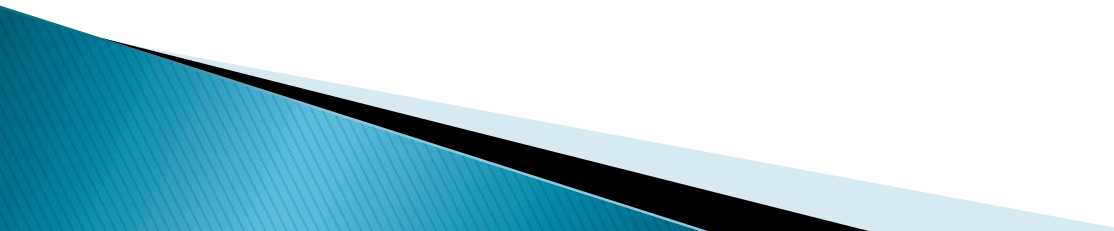
- ▶ Each router (switch) makes a LOCAL decision to forward the packet towards B



Connectionless Forwarding

- ▶ This is termed *destination-based connectionless forwarding*
- ▶ How does each router know the *correct* local forwarding decision for any possible destination address?
 - Through knowledge of the *topology state* of the network
 - This knowledge is maintained by a *routing protocol*

Routing Protocols

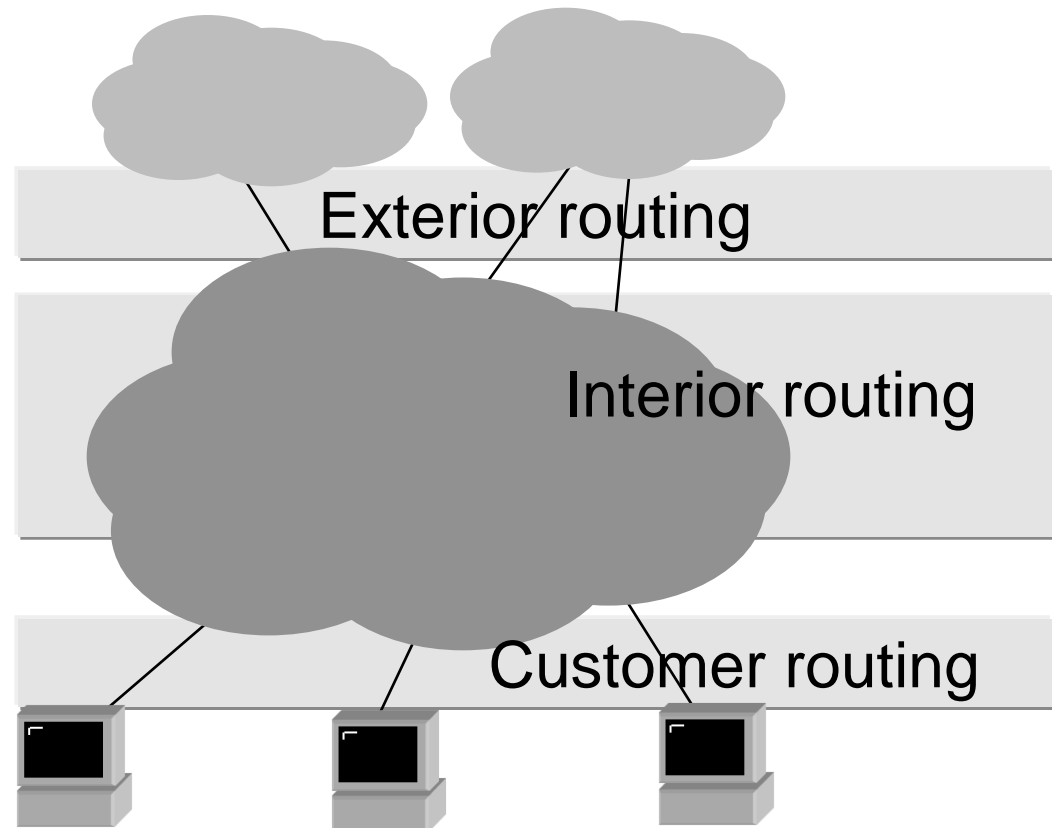
- ▶ Distribute the knowledge of the current *topology state* of the network to all routers
 - ▶ This knowledge is used by each router to generate a *forwarding table*, which contains the local switching decision for each known destination address
- 

Routing Protocols

- ▶ correct operation of the routing state of a network is essential for the management of a *quality* network service
 - accuracy of the routing information
 - dynamic adjustment of the routing information
 - matching aggregate traffic flow to network capacity

ISP Routing Tasks

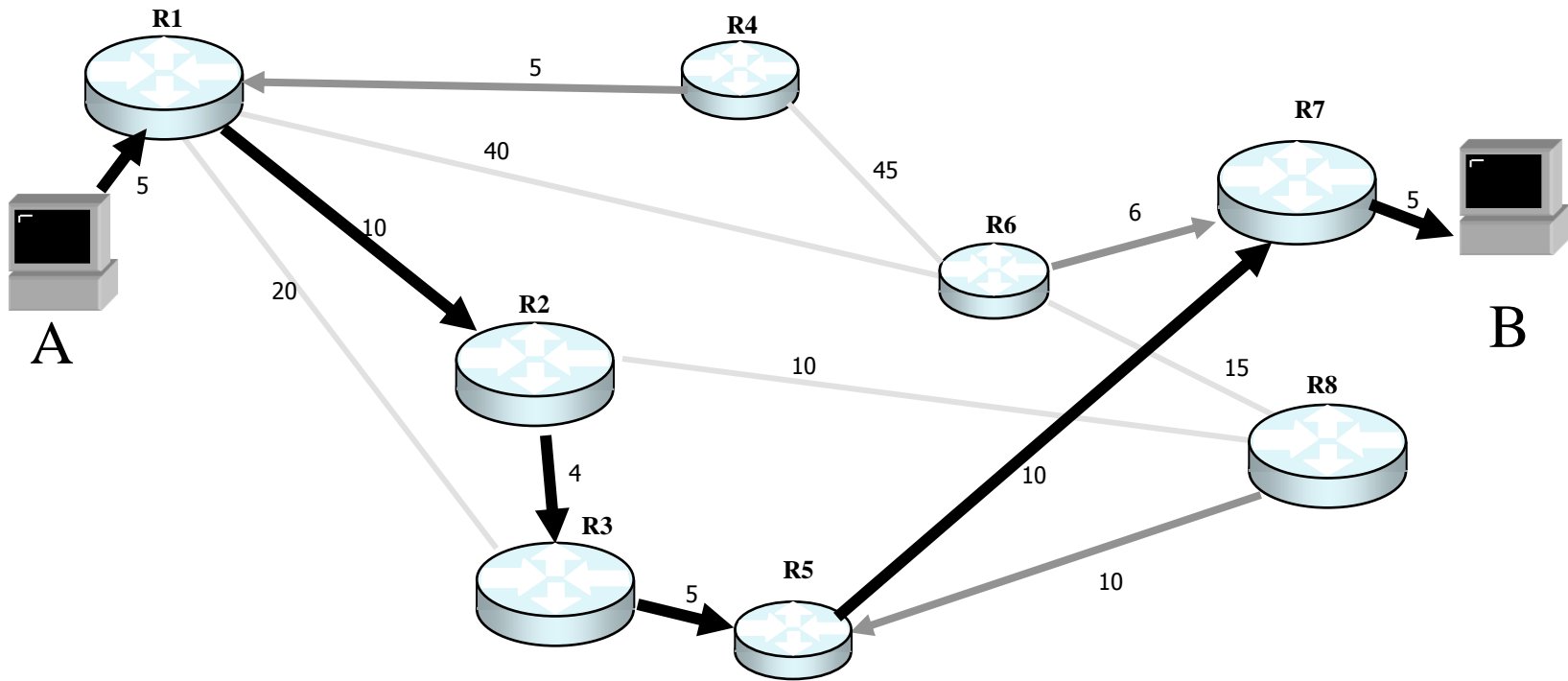
- ▶ customers
- ▶ internal
- ▶ peer / upstream



Interior Routing

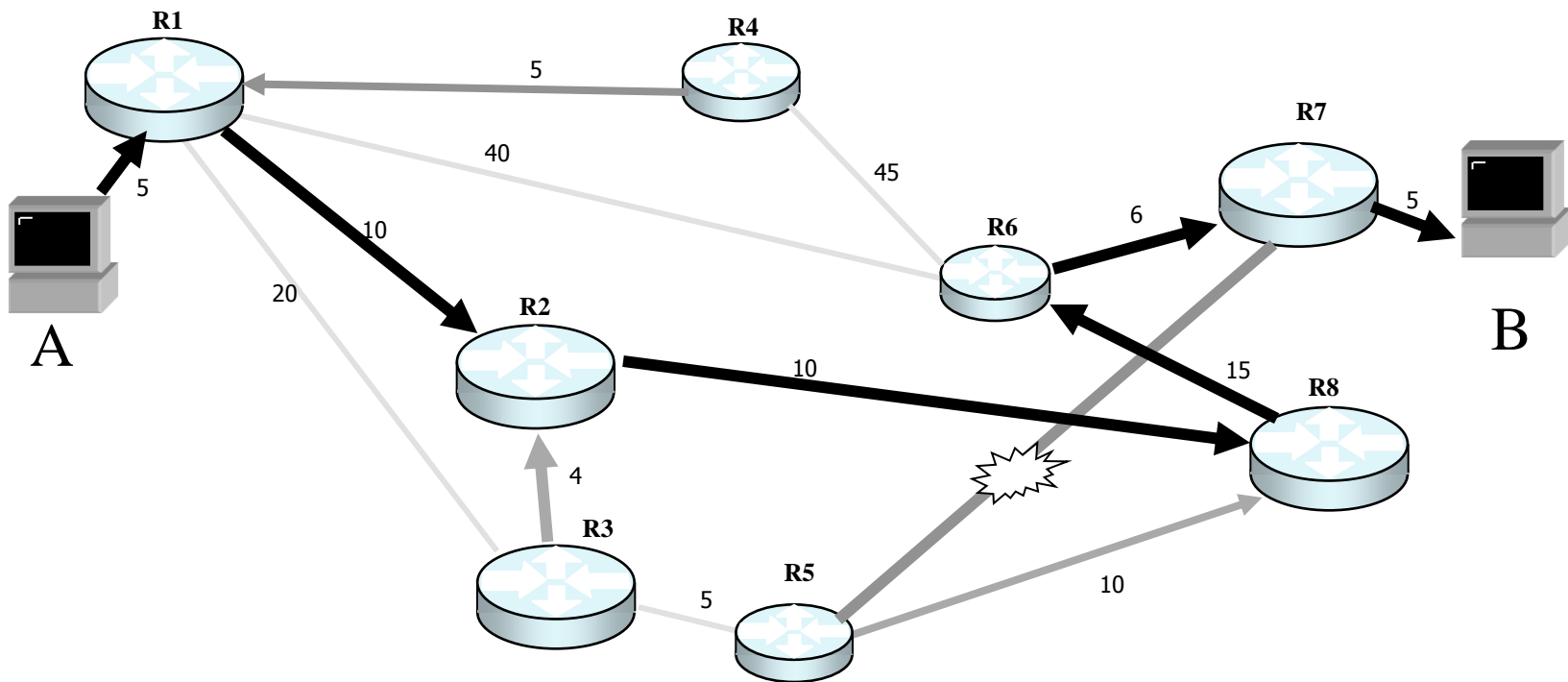
- ▶ discovers the *topology* of a network through the operation of a *distributed routing protocol*

Path Selection



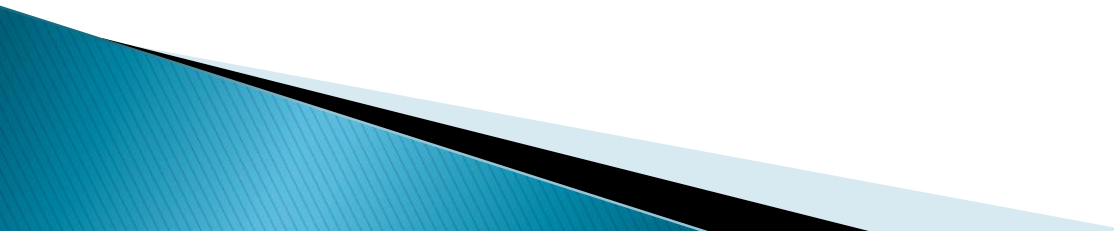
Minimum cost from A to B is 39 units

Dynamic Path Adjustment



If R5 – R7 breaks, minimum cost path from A to B is
Now 46 units

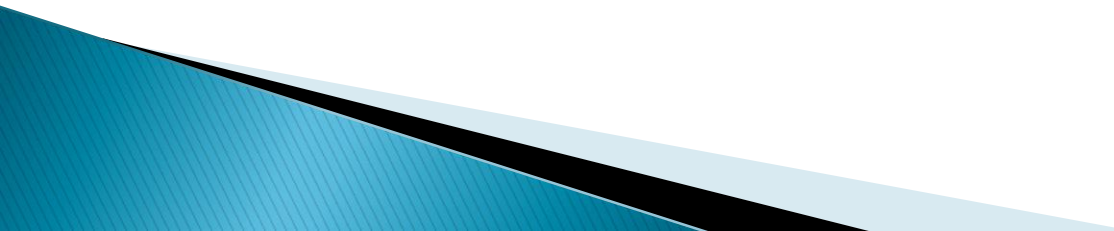
Interior Routing Protocols

- ▶ describe the current network topology
 - ▶ Routing protocols distribute **how** to reach address prefix groups
 - ▶ Routing protocols function through either
 - distributed computing model (distance vector)
 - parallel computing model (link state)
- 

Routing Protocols

- ▶ Distance Vector Routing Protocols
 - Each node sends its routing table (dest, distance) to all neighbors every 30 seconds
 - Lower distances are updated with the neighbor as next hop
 - cannot scale
 - cannot resolve routing loops quickly
 - RIP is the main offender

Routing Protocols

- ▶ Link State Routing Protocols
 - Each link, the connected nodes and the metric is flooded to all routers
 - Each link up/down status change is incrementally flooded
 - Each router re-computes the routing table in parallel using the common link state database
 - OSPF is the main protocol in use today
- 

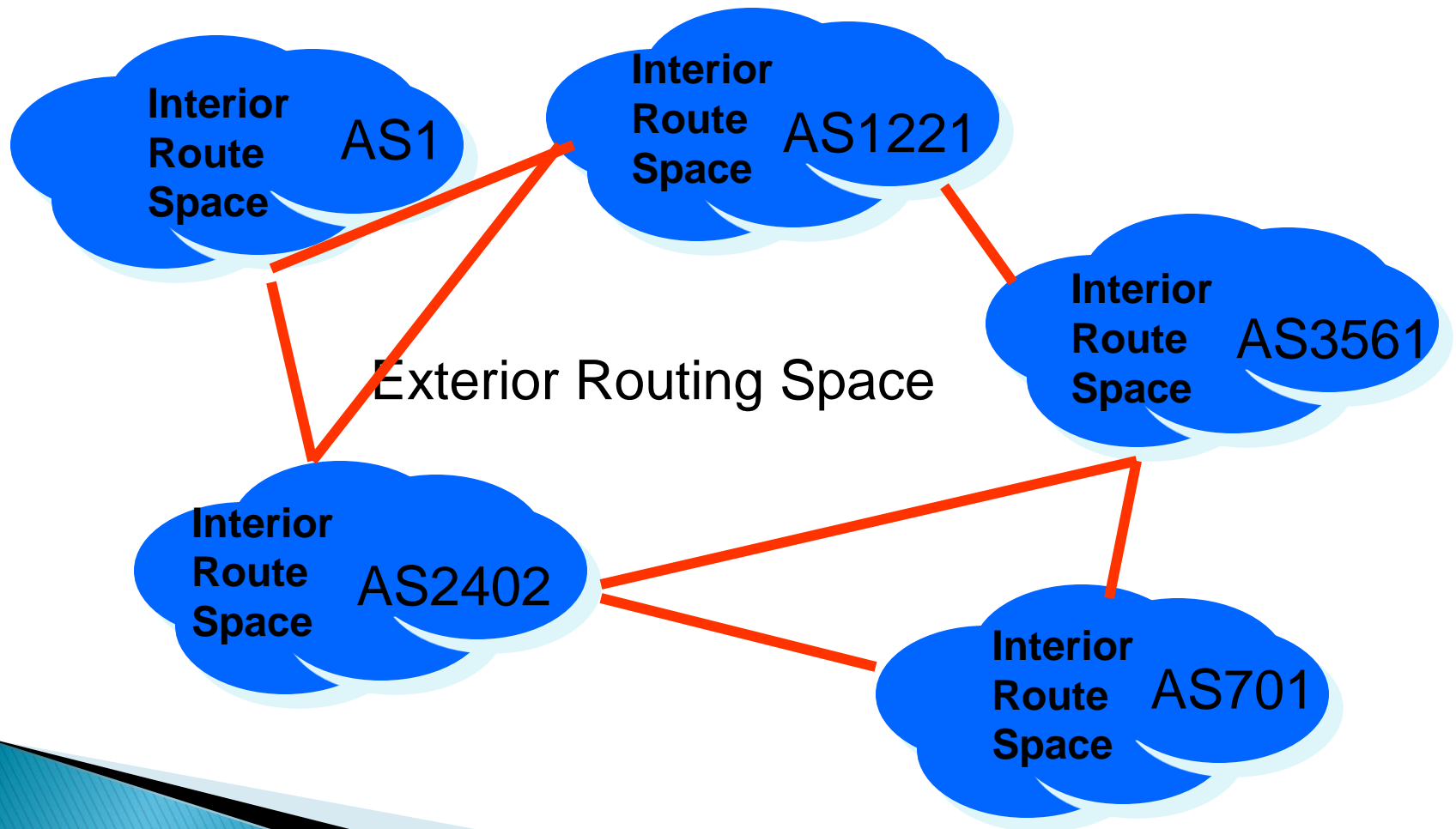
Suggestions

- ▶ Just engineering a physical link does not ensure that traffic will flow
 - some system somewhere must provide routing information about how to reach the newly connected network
- ▶ Installing backup circuits is easy, making the routing work may not be

Suggestions

- ▶ need a clear understanding of how the client networks want their traffic to flow before you can start making routing configuration changes

Interior and Exterior Routing Protocols



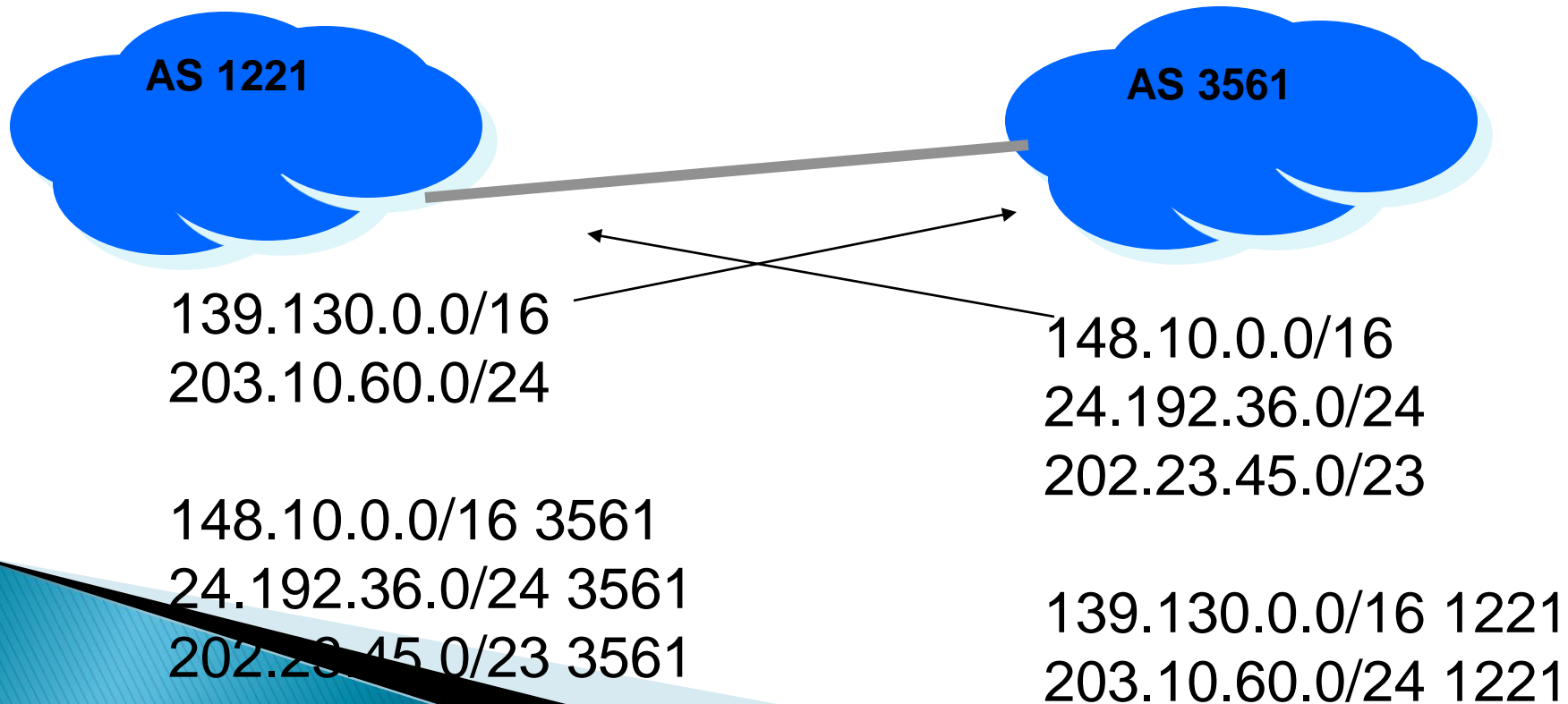
Exterior Routing Protocols

- ▶ You tell me all the address prefixes you can reach, but don't tell me the path you use to get there
 - I'll tell you the same
- ▶ If anything changes, please let me know
- ▶ If you tell me an address I'll send you traffic destined to that address.
 - If I tell you an address I will accept traffic destined to that address

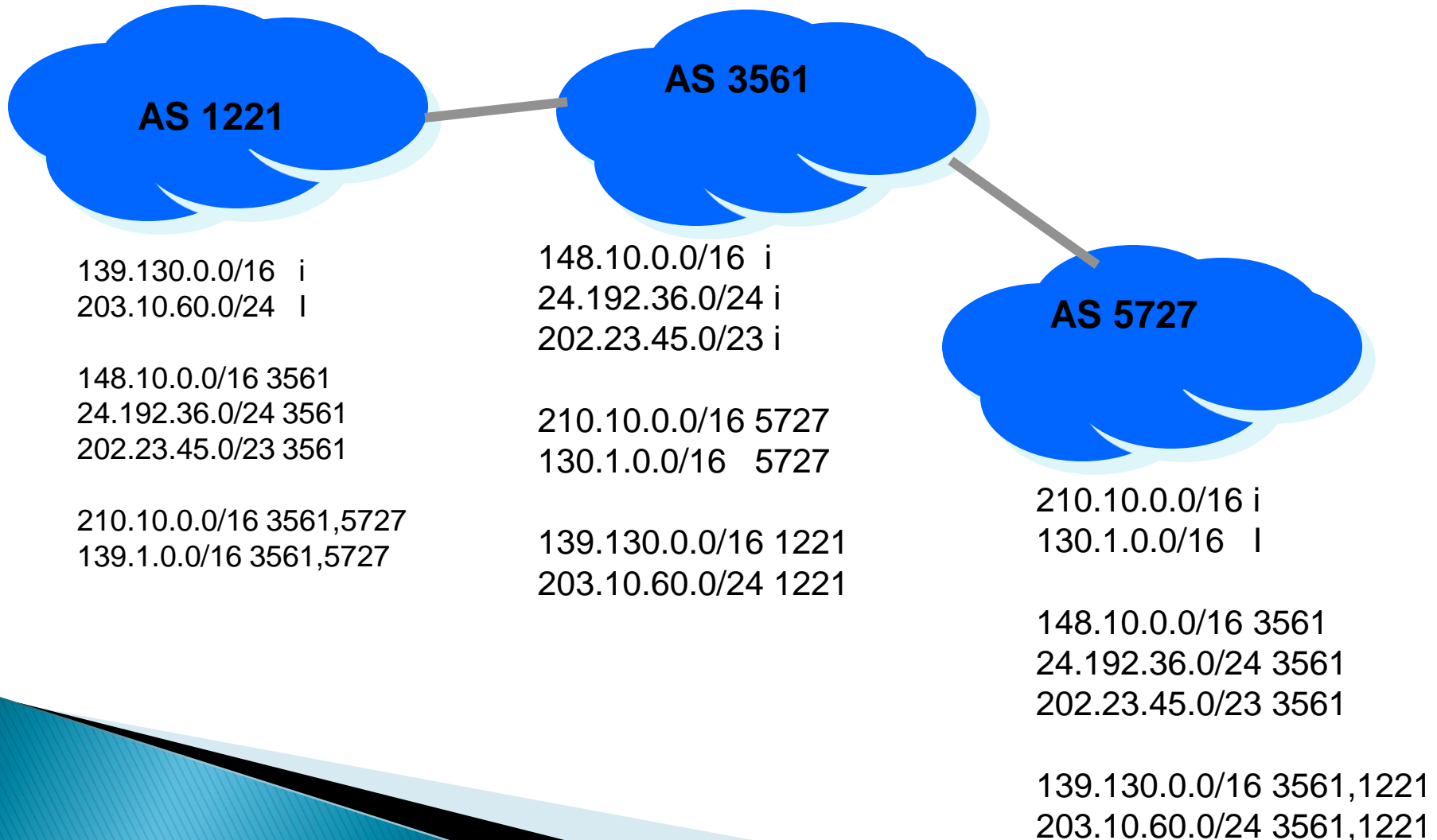
Exterior Routing Protocols

- ▶ Border Gateway Protocol version 4 (BGP4)
- ▶ Each interior route collection is described by an Autonomous System (AS) number
- ▶ Internal topology is hidden
- ▶ Routes are announced with associated AS value
 - 139.130.0.0/16 + AS 1221

BGP example

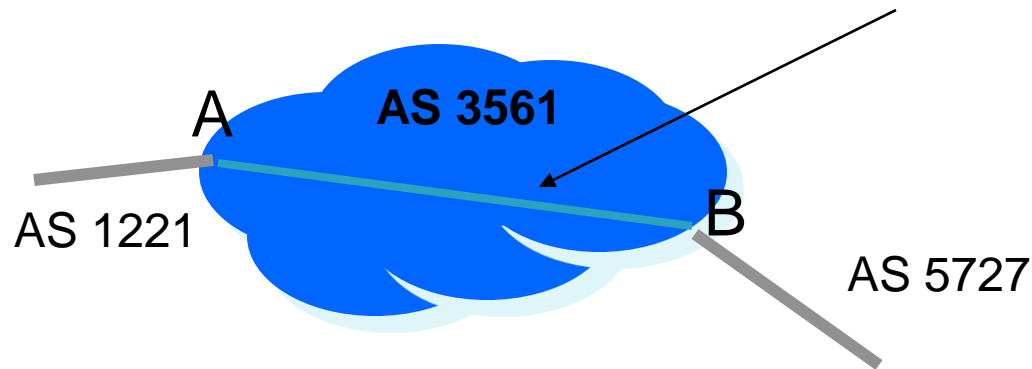


BGP Example of TRANSIT



Exterior Routing Protocols

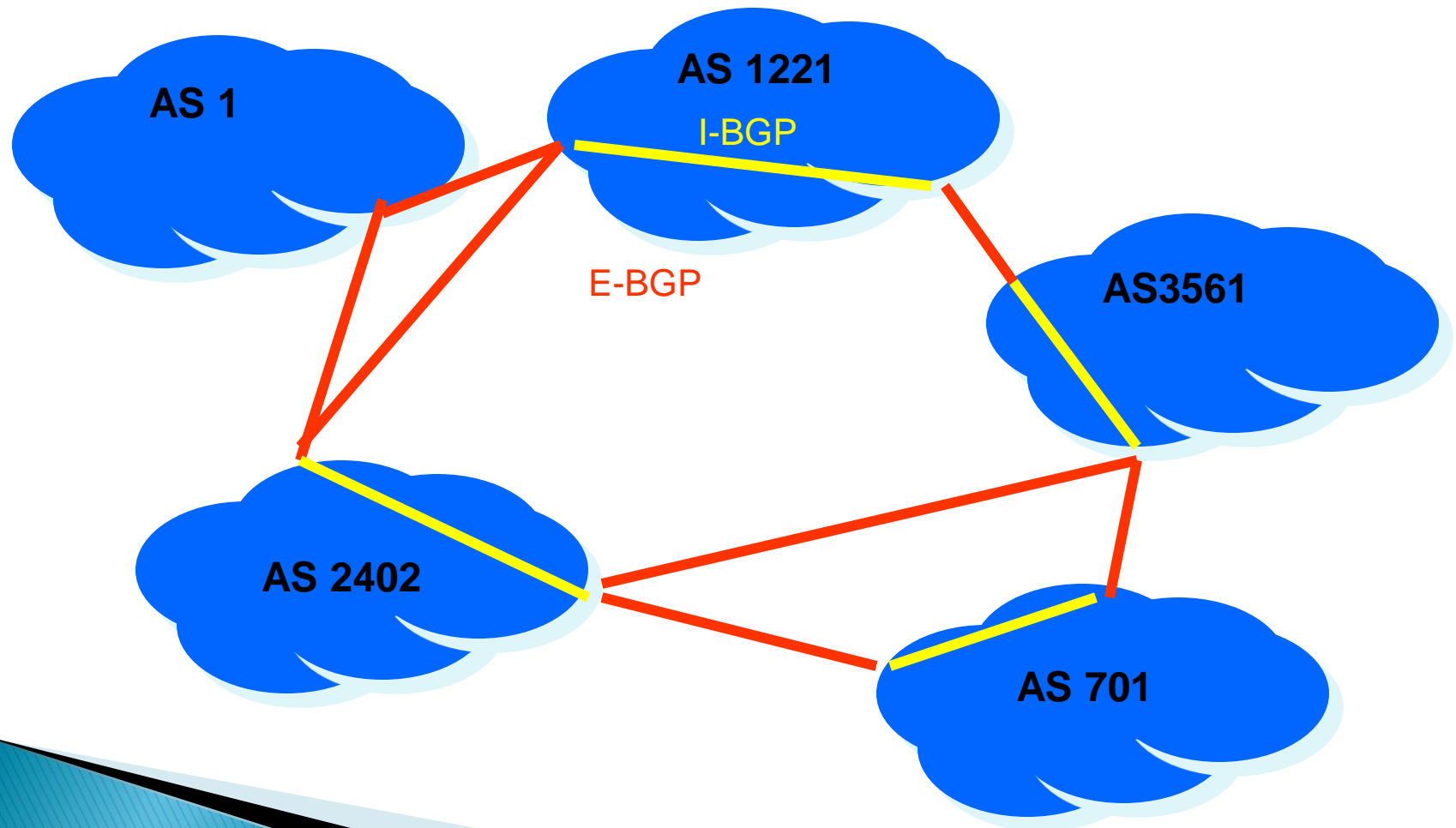
- ▶ Internal transit paths use I-BGP



Q: How does router A tell router B about AS1221 addresses?

A: Router A sets up an INTERIOR BGP session with router B

Exterior Routing Protocols



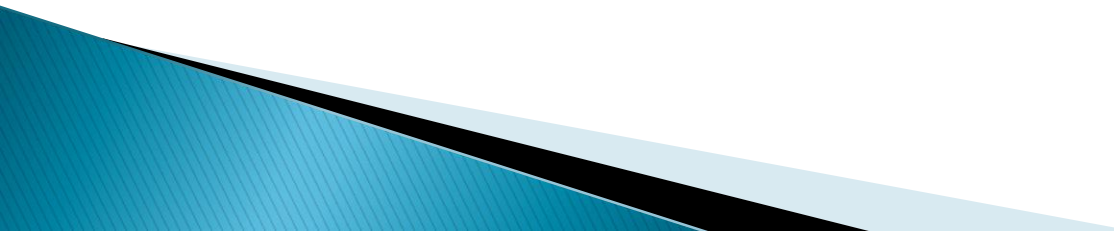
Exterior Routing Protocols

- ▶ Normally chose minimal AS path length

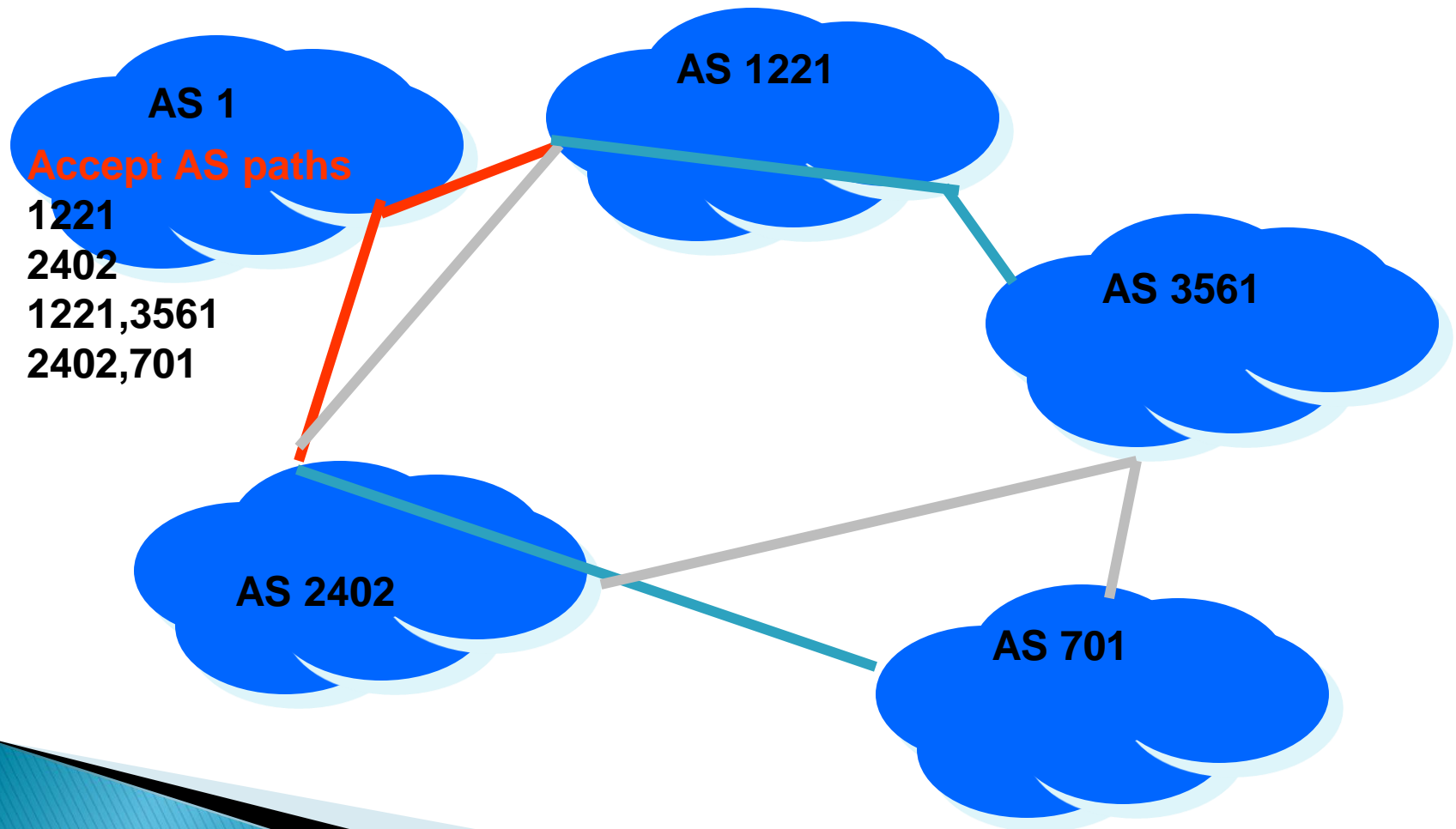
203.10.60.0/24 701,3561,1221
—————→ 203.10.60.0/24 5727,1221

Selected path is via peer session to AS 5727 as this is 1 AS shorter than the other path

Exterior POLICY

- ▶ How can I share the traffic load between 2 or more exterior providers?
 - ▶ How can I create a backup link to support my main exterior link?
 - ▶ You can bias minimal path selection by **AS path filter lists** or **community attributes** or **local preferences**
- 

Exterior Routing Protocols plus Policy



Exterior Routing Protocols plus Policy

- ▶ policy settings control
 - what you advertise to your immediate peers
 - What you accept from your immediate peers
 - What transits you will accept (send traffic)
- ▶ **you cannot control**
 - transit path of received traffic
 - symmetry of transit policy